




# Transportation Problems

- Transportation is considered as a “special case” of LP
- Reasons?
  - it can be formulated using LP technique so is its solution



- Here, we attempt to firstly define what are them and then studying their solution methods:  (to p3)

# Transportation Problem

- We have seen a sample of transportation problem on slide 29 in lecture 2  (to p4)
- Here, we study its alternative solution method
- Consider the following transportation tableau  (to p6)

# Review of Transportation Problem

Warehouse supply of televisions sets:

1- Cincinnati	300
2- Atlanta	200
3- Pittsburgh	200
total	700

Retail store demand for television sets:

A - New York	150
B - Dallas	250
C - Detroit	200
total	600

From Warehouse	To Store		
	A	B	C
1	\$16	\$18	\$11
2	14	12	13
3	13	15	17

LP formulation



(to p5)

# A Transportation Example (2 of 3)

## Model Summary and Computer Solution with Excel

Minimize  $Z = \$16x_{1A} + 18x_{1B} + 11x_{1C} + 14x_{2A} + 12x_{2B} + 13x_{2C} + 13x_{3A} + 15x_{3B} + 17x_{3C}$

subject to

$$x_{1A} + x_{1B} + x_{1C} \leq 300$$

$$x_{2A} + x_{2B} + x_{2C} \leq 200$$

$$x_{3A} + x_{3B} + x_{3C} \leq 200$$

$$x_{1A} + x_{2A} + x_{3A} = 150$$

$$x_{1B} + x_{2B} + x_{3B} = 250$$

$$x_{1C} + x_{2C} + x_{3C} = 200$$

$$x_{ij} \geq 0$$

Objective function

	A	B	C	D	E	F	G	H
1	A Transportation Example							
2								
3				Store			TV Sets	
4	Warehouse	New York	Dallas	Detroit	Supply	Shipped		
5	Cincinnati	0	0	200	300	200		
6	Atlanta	0	200	0	200	200		
7	Pittsburgh	150	50	0	200	200		
8	Demand	150	250	200				
9	TV Sets Shipped	150	250	200				
10	Cost = \$	7300						
11								

=C5+C6+C7

=C5+D5+E5

Exhibit 4.15


(to p3)







# Transportation Tableau

Table B-1  
The Transportation Tableau





From \ To	A	B	C	Supply
1	6	8	10	150
2	7	11	11	175
3	4	5	12	275
Demand	200	100	300	600

- We know how to formulate it using LP technique
  - Refer to lecture 2 note
- Here, we study its solution by firstly attempting to determine its initial tableau  (to p7)
  - Just like the first simplex tableau!

# Solution to a transportation problem

- Initial tableau  (to p8)
- Optimal solution  (to p23)
- Important Notes  (to p43)
- Tutorials  (to p53)

# initial tableau

- Three different ways:
  - Northwest corner method  (to p9)
  - The Minimum cell cost method  (to p12)
  - Vogel's approximation method (VAM)  (to p16)
- Now, are these initial tableaus given us an Optimal solution?  (to p7)



# Northwest corner method

## Steps:

1. assign largest possible allocation to the cell in the upper left-hand corner of the tableau
2. Repeat step 1 until all allocations have been assigned
3. Stop. Initial tableau is obtained

Example  (to p10)

# Northeast corner

Step 1

Step 2

Max (150,200)

Table B-1  
The Transportation Tableau

From \ To	A	B	C	Supply
1	6 150	8 --	10 --	150
2	7	11	11	175
3	4	5	12	275
Demand	200	100	300	600

0

150

Table B-1  
The Transportation Tableau

From \ To	A	B	C	Supply
1	6 150	8 --	10 --	150
2	7 50	11	11	175
3	4 --	5	12	275
Demand	200	100	300	600

0

125

0

(to p11)



# Initial tableau of NW corner method

- Repeat the above steps, we have the following tableau.
- Stop. Since all allocated have been assigned

**Table B-2**  
The Initial NW Corner Solution

From \ To	A	B	C	Supply
1	6 150	8	10	150
2	7 50	11 100	11 25	175
3	4	5	12 275	275
<b>Demand</b>	200	100	300	600

Ensure that all columns and rows added up to its respective totals.

(to p8)



# The Minimum cell cost method

Here, we use the following steps:

Steps:

Step 1 Find the cell that has the least cost

Step 2: Assign as much as allocation to this cell

Step 3: Block those cells that cannot be allocated

Step 4: Repeat above steps until all allocation have been assigned.

Example:  (to p13)

Step 1 Find the cell that has the least cost  
 Step 2: Assign as much as allocation to this cell

Step 1:

Table B-1  
The Transportation Tableau

From \ To	A	B	C	Supply
1	6	8	10	150
2	7	11	11	175
3	4	5	12	275
Demand	200	100	300	600

Step 2:

Table B-1  
The Transportation Tableau

From \ To	A	B	C	Supply
1	6	8	10	150
2	7	11	11	175
3	4	5	12	275
Demand	200	100	300	600



The min cost, so allocate as much resource as possible here

Step 3



(to p14)

Step 3: Block those cells that cannot be allocated

Step 4: Repeat above steps until all allocation have been assigned.

Step 3:

Table B-1  
The Transportation Tableau

From \ To	A	B	C	Supply
1	--	6	8	150
2	--	7	11	175
3	200	4	5	<del>275</del> 75
Demand	<del>200</del>	100	300	600

0

Second iteration, step 4

Table B-4  
The Second Minimum Cell Cost Allocation

From \ To	A	B	C	Supply
1		6	8	150
2		7	11	175
3	200	4	5	75
Demand	200	100	300	600

(to p15)

# The initial solution

**Table B-5**  
The Initial Solution

From \ To	A	B	C	Supply
1	6	8	10	150
2	7	11	11	175
3	4	5	12	275
Demand	200	100	300	600

- Stop. The above tableau is an initial tableau because all allocations have been assigned



# Vogel's approximation method

## Operational steps:

- Step 1: for each column and row, determine its penalty cost by subtracting their two of their least cost
- Step 2: select row/column that has the highest penalty cost in step 1
- Step 3: assign as much as allocation to the selected row/column that has the least cost
- Step 4: Block those cells that cannot be further allocated
- Step 5: Repeat above steps until all allocations have been assigned

Example  (to p17)



# subtracting their two of their least cost

Step 1

**Table B-6**  
The VAM Penalty Costs

From \ To	A	B	C	Supply
1	6	8	10	150
2	7	11	11	175
3	4	5	12	275
Demand	200	100	300	600

2 (8-6)

4 (11-7)

1 (5-4)

2 (6-4)      3 (8-5)      1 (11-10)

(to p18)



# Steps 2 & 3

**Table B-6**  
The VAM Penalty Costs

From \ To	A	B	C	Supply
1	6	8	10	150
2	7	11	11	175
3	4	5	12	275
Demand	200	100	300	600

**Step 2:**  
Highest penalty cost

2  
4  
1

(to p19)



**Step 3:** this has the least cost

# Step 4

**Table B-7**  
The Initial VAM Allocation

From \ To	A	B	C	Supply	
1	6	8	10	150	2
2	7	11	11	175	1
3	4	5	12	275	
Demand	200	100	300	600	
	2	3	1		

(to p20)





# 3<sup>rd</sup> Iteration of VAM

**Table B-9**  
The Third VAM Allocation

From \ To	A	B	C	Supply
1	---	---	---	150
2	175	---	---	175
3	25	100	---	275
Demand	200	100	300	600

2

(to p22)



# Initial tableau for VAM

**Table B-10**  
The Initial VAM Solution

From \ To	A	B	C	Supply
1	6	8	10	150
2	7	11	11	175
3	4	5	12	275
Demand	200	100	300	600

(to p8)



# Optimal solution?

Initial solution from:

Northeast cost, total cost = \$5,925




The min cost, total cost = \$4,550

VAM, total cost = \$5,125

(note: here, we are not saying the second one always better!)

It shows that the second one has the min cost, but is it the optimal solution?  (to p24)

# Solution methods

- We need a method, like the simplex method, to check and obtain the optimal solution
- Two methods:
  1. Stepping-stone method  (to p25)  (to p30)
  2. Modified distributed method (MODI) 





# Stepping-stone method

Let consider the following initial tableau from the Min Cost algorithm

**Table B-11**  
The Minimum Cell Cost Solution

From \ To	A	B	C	Supply
1	6	8	10	150
2	7	11	11	175
3	4	5	12	275
Demand	200	100	300	600

There are Non-basic variables

These are basic variables

(to p26)

Question: How can we introduce a non-basic variable into basic variable?



# Introduce

a non-basic variable into basic variables

- Here, we can select any non-basic variable as an entry and then using the “+ and –” steps to form a closed loop as follows:

**Table B-11**  
The Minimum Cell Cost Solution

From \ To	A	B	C	Supply
1	6	8	10	
2	7	11	11	175
3	4	5	12	275
Demand	200	100	300	600

let consider this non basic variable

1

Then we have



(to p27)

# Stepping stone

**Table B-11**  
The Minimum Cell  
Cost Solution

From \ To	A	B	C	Supply
1	+ 6	- 8	10	150
2	7	11	11	175
3	4	+ 5	12	275
Demand	200	100	300	600

Diagram annotations: A dotted path starts at cell (1, B) with a value of 25. It moves left to cell (1, A) with a value of 6, then down to cell (3, A) with a value of 4, and finally right to cell (3, B) with a value of 5. A '+' sign is placed in cell (1, A) and a '-' sign is placed in cell (1, B). Dotted arrows indicate the direction of the path: left from (1, B) to (1, A), down from (1, A) to (3, A), and right from (3, A) to (3, B).

(to p28)



The above saying that, we add min value of all -ve cells into cell that has "+" sign, and subtracts the same value to the "-ve" cells

Thus, max -ve is  $\min(200, 25) = 25$ , and we add 25 to cell A1 and A3, and subtract it from B1 and A3

# Stepping stone

**Table B-19**  
The Second Iteration of the  
Stepping-Stone Method

From \ To	A	B	C	Supply
1	25	8	10	150
2	7	11	11	175
3	4	5	12	275
Demand	200	100	300	600

The above tableaus give min cost =  $25 \cdot 6 + 120 \cdot 10 + 175 \cdot 11$   
 $175 \cdot 4 + 100 \cdot 5 = \$4525$

We can repeat this process to all possible non-basic cells in that above tableau until one has the min cost! → NOT a Good solution method

(to p29)



# Getting optimal solution

- In such, we introducing the next algorithm called Modified Distribution (MODI)

(to p24)



# Modified distributed method (MODI)

- It is a modified version of stepping stone method
- MODI has two important elements:
  1. It determines if a tableau is the optimal one
  2. It tells you which non-basic variable should be firstly considered as an entry variable
  3. It makes use of stepping-stone to get its answer of next iteration

– How it works?



(to p31)

# Procedure (MODI)

Step 0: let  $u_i$ ,  $v_j$ ,  $c_{ij}$  variables represent rows, columns, and cost in the transportation tableau, respectively

Step 1: (a) form a set of equations that uses to represent all basic variables

$$u_i + v_j = c_{ij}$$

(b) solve variables by assign one variable = 0

Step 2: (a) form a set of equations use to represent non-basic variable (or empty cell) as such

$$c_{ij} - u_i - v_j = k_{ij}$$

(b) solve variables by using step 1b information

Step 3: Select the cell that has the most -ve value in 2b

Step 4: Use stepping-stone method to allocate resource to cell in step 3

Step 5: Repeat the above steps until all cells in 2a has no negative value

Example



(to p32)

(to p24)



# MODI

Consider to this initial tableau:

**Table B-11**  
The Minimum Cell Cost Solution

From \ To	A	B	C	Supply
1	6	8	10	150
2	7	11	11	175
3	4	5	12	275
Demand	200	100	300	600

Step 0: let  $u_i$ ,  $v_j$ ,  $c_{ij}$  variables represent rows, columns, and cost in the transportation tableau, respectively



(to p33)



# Step 0

	$v_j$	$v_A =$	$v_B =$	$v_C =$	
$u_i$	From \ To	A	B	C	Supply
$u_1 =$	1	6	8	10	150
$u_2 =$	2	7	11	11	175
$u_3 =$	3	4	5	12	275
	Demand	200	100	300	600

$C_{3A}$

Step 1: (a) form a set of equations that uses to represent all basic variables (to p34)

$$u_i + v_j = c_{ij}$$



# $u_i + v_j = c_{ij}$

$$u_i + v_j = c_{ij}$$

The value  $c_{ij}$  is the unit transportation cost for cell  $ij$ . For example, the formula for cell 1B is

$$u_1 + v_B = c_{1B}$$

and, since  $c_{1B} = 8$ ,

$$u_1 + v_B = 8$$

The formulas for the remaining cells that presently contain allocations are

$$x_{1C}: u_1 + v_C = 10$$

$$x_{2C}: u_2 + v_C = 11$$

$$x_{3A}: u_3 + v_A = 4$$

$$x_{3B}: u_3 + v_B = 5$$

	$v_j$	$v_A =$	$v_B =$	$v_C =$	
	<b>To</b>				
$u_i$	<b>From</b>	A	B	C	<b>Supply</b>
$u_1 =$	1	6	8	10	150
$u_2 =$	2	7	11	11	175
$u_3 =$	3	4	5	12	275
	<b>Demand</b>	200	100	300	600

(to p35)

(b) solve variables by assign  
one variable = 0



# Set one variable = 0

Now there are five equations with six unknowns. To solve these equations, it is necessary to assign only one of the unknowns a value of zero. Thus, if we let  $u_1 = 0$ , we can solve for all remaining  $u_i$  and  $v_j$  values.

Because we added a non-basic variable

$$x_{1B}: u_1 + v_B = 8$$

$$0 + v_B = 8$$

$$v_B = 8$$

$$x_{1C}: u_1 + v_C = 10$$

$$0 + v_C = 10$$

$$v_C = 10$$

$$x_{2C}: u_2 + v_C = 11$$

$$u_2 + 10 = 11$$

$$u_2 = 1$$

$$x_{3B}: u_3 + v_B = 5$$

$$u_3 + 8 = 5$$

$$u_3 = -3$$

$$x_{3A}: u_3 + v_A = 4$$

$$-3 + v_A = 4$$

$$v_A = 7$$

(to p36)

Step2: (a & b)



# Step2: (a & b)

	$v_j$	$v_A = 7$	$v_B = 8$	$v_C = 10$	
$u_i$	From \ To	A	B	C	Supply
$u_1 = 0$	1	6	25	125	150
$u_2 = 1$	2	7	11	175	175
$u_3 = -3$	3	4	5	12	275
	<b>Demand</b>	200	100	300	600

Next, we use the following formula to evaluate all *empty cells*:

$$c_{ij} - u_i - v_j = k_{ij}$$

where  $k_{ij}$  equals the cost increase or decrease that would occur by allocating to a cell.

For the *empty cells* in Table B-26, the formula yields the following values:

$$\begin{aligned} x_{1A}: k_{1A} &= c_{1A} - u_1 - v_A = 6 - 0 - 7 = -1 \\ x_{2A}: k_{2A} &= c_{2A} - u_2 - v_A = 7 - 1 - 7 = -1 \\ x_{2B}: k_{2B} &= c_{2B} - u_2 - v_B = 11 - 1 - 8 = +2 \\ x_{3C}: k_{3C} &= c_{3C} - u_3 - v_C = 12 - (-3) - 10 = +5 \end{aligned}$$

Note this may look difficult and complicated, however, we can add these V=values into the above tableau as well



(to p37)

# Step2: (a & b), alternative

		$v_j$	$v_A = 7$	$v_B = 8$	$v_C = 10$	
$u_i$	From \ To	A	B	C	Supply	
$u_1 = 0$	1	-1 6	25 8	125 10	150	
$u_2 = 1$	2	-1 7	+2 11	175 11	175	
$u_3 = -3$	3	200 4	75 5	+5 12	275	
	Demand	200	100	300	600	

Next, we use the following formula to evaluate all *empty cells*:

$$c_{ij} - u_i - v_j = k_{ij}$$

6-0-7

(to p38)

Step 3: Select the cell that has the most -ve value in 2b



# Step3

	$v_j$	$v_A = 7$	$v_B = 8$	$v_C = 10$	
$u_i$	From \ To	A	B	C	Supply
$u_1 = 0$	1	-1 6	25 8	125 10	150
$u_2 = 1$	2	-1 7	+2 11	175 11	175
$u_3 = -3$	3	4	5	+5 12	275
	Demand	200	100	300	600

Next, we use the following formula to evaluate all *empty cells*:

$$c_{ij} - u_i - v_j = k_{ij}$$

Select either one, (Why?)

These cells mean, introduce it will reduce the min z to -1 cost unit

(to p39)



# Step 4: Use stepping-stone method

**Table B-11**  
The Minimum Cell  
Cost Solution

From \ To	A	B	C	Supply
1	+ 6	- 8	10	150
2	7	11	11	175
3	4	+ 5	12	275
Demand	200	100	300	600

Diagram illustrating the stepping-stone method. A path is shown starting from cell (1, A) with a '+' sign, moving left to cell (1, B) with a '-' sign, then down to cell (2, B) with a '+' sign, then down to cell (3, B) with a '-' sign, and finally right to cell (3, A) with a '+' sign. A value of 25 is shown in the center of the path, indicating the amount to be shifted.

From here we have ....



(to p40)

# Step 4: Use stepping-stone method

**Table B-19**  
The Second Iteration of the  
Stepping-Stone Method

From \ To	A	B	C	Supply
1	25		125	150
2			175	175
3	175	100		275
Demand	200	100	300	600

(to p41)

Step 5: we repeat steps 1-4 again for the above tableau, we have






# Step 5

**Table B-27**  
The Second Iteration of the  
MODI Solution Method

	$v_j$	$v_A =$	$v_B =$	$v_C =$	
$u_i$	From \ To	A	B	C	Supply
$u_1 =$	1	25	125	150	150
$u_2 =$	2	175	175	175	175
$u_3 =$	3	275	275	275	275
	Demand	200	100	300	600



$$\begin{aligned}
 x_{1A}: \quad & u_1 + v_A = 6 \\
 & 0 + v_A = 6 \\
 & v_A = 6 \\
 x_{1C}: \quad & u_1 + v_C = 10 \\
 & 0 + v_C = 10 \\
 & v_C = 10 \\
 x_{2C}: \quad & u_2 + v_C = 11 \\
 & u_2 + 10 = 11 \\
 & u_2 = 1 \\
 x_{3A}: \quad & u_3 + v_A = 4 \\
 & u_3 + 6 = 4 \\
 & u_3 = -2 \\
 x_{3B}: \quad & u_3 + v_B = 5 \\
 & -2 + v_B = 5 \\
 & v_B = 7
 \end{aligned}$$

(to p42) 



# Step 5 cont

**Table B-28**  
The New  $u_i$  and  $v_j$  Values for  
the Second Iteration

	$v_j$	$v_A = 6$	$v_B = 7$	$v_C = 10$	
$u_i$	To	A	B	C	Supply
$u_1 = 0$	From	6	8	10	
	1	25		125	150
$u_2 = 1$	2	7	11	11	175
	3	4	5	12	
$u_3 = -2$		175	100		275
	Demand	200	100	300	600

The cost changes for the empty cells are now computed using the formula  
 $c_{ij} - u_i - v_j = k_{ij}$

$$\begin{aligned}
 x_{1A}: k_{1B} &= c_{1B} - u_1 - v_B = 8 - 0 - 7 = +1 \\
 x_{2A}: k_{2A} &= c_{2A} - u_2 - v_A = 7 - 1 - 6 = 0 \\
 x_{2B}: k_{2B} &= c_{2B} - u_2 - v_B = 11 - 1 - 7 = +3 \\
 x_{3C}: k_{3C} &= c_{3C} - u_3 - v_C = 12 - (-2) - 10 = +4
 \end{aligned}$$



All positives  
**STOP**


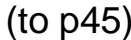



Because none of these values is negative, the solution shown in Table B-28 is optimal. However, as in the stepping-stone method, cell 2A with a zero cost change indicates a multiple optimal solution.

(to p31)



# Important Notes

- When start solving a transportation problem using algorithm, we need to ensure the following:

1. Alternative solution  (to p44)  (to p45)
2. Total demand  $\neq$  total supply 
3. Degeneracy  (to p48)
4. others  (to p52)



# Alternative solution

- When on the following  $k = 0$

$$c_{ij} - u_i - v_j = k_{ij}$$

Why?



# Total demand $\neq$ total supply

From \ To	A	B	C	Supply
1	6	8	10	150
2	7	11	11	175
3	4	5	12	275
Demand	200	100	350	

Note that, total demand=650, and total supply = 600

How to solve it?

(to p46)

We need to add a dummy row and assign 0 cost to each cell as such ..



# Dd≠ss

From \ To	A	B	C	Supply
1	6	8	10	150
2	7	11	11	175
3	4	5	12	275
Dummy	0	0	0	50
Demand	200	100	350	650

Extra row, since Demand > supply

Other example .....



(to p47)

# Dd≠ss

**Table B-30**  
An Unbalanced Model  
(Supply > Demand)

From \ To	A	B	C	Dummy	Supply
1	6	8	10	0	150
2	7	11	11	0	175
3	4	5	12	0	375
Demand	200	100	300	100	700

Extra column is added

(to p43)



# Degeneracy

In all the tableaus showing a solution to the wheat transportation problem, the following condition was met.

$$m \text{ rows} + n \text{ columns} - 1 = \text{the number of cells with allocations}$$

Example .....



(to p49)

ie, basic variables  
in the tableau



# Degeneracy


**Table B-11**  
The Minimum Cell  
Cost Solution

From \ To	A	B	C	Supply
1	6	8	10	150
2	7	11	11	175
3	4	5	12	275
Demand	200	100	300	600

$m$  rows +  $n$  column - 1 = the number of cells with allocations

$3 + 3 - 1 = 5$  (five basic variables, and above has 5 as well!)

It satisfied.

If failed? ..... considering .....  (to p50)

# Degeneracy

**Table B-31**  
The Minimum Cell Cost Initial Solution

From \ To	A	B	C	Supply
1	6	8	10	150
2	7	11	11	250
3	4	5	12	200
<b>Demand</b>	200	100	300	600

The tableau shown in Table B-31 does not meet the condition

$$m + n - 1 = \text{the number of cells with allocations}$$

$$3 + 3 - 1 = 5 \text{ cells}$$

(note above has only 4 basic variable only!)

If not matched, then we select a non-basic variable with least cheapest cost and considered it as a new basic variable with assigned 0 allocation to it

(to p51)



# Degeneracy

**Table B-32**  
The Initial Solution

From \ To	A	B	C	Supply
1	0	100	50	150
2			250	250
3	200			200
Demand	200	100	300	600

Added this

Note: we pick this over others because it has the least cost for the Min Z problem!

# others

## 1. When one route cannot be used

- Assign a big M cost to its cell

If 10 changed to  
Cannot delivered  
Then we assigned  
M value here

Table B-32  
The Initial Solution

From \ To	A	B	C	Supply
1	6 0	8 100	10 50	150
2	7	11	11 250	250
3	4 200	5	12	200
Demand	200	100	300	600

(to p43)



# Tutorials

- Module B
  - 1, 5, 8, 13, 21, 34
  - these questions are attached in the following slides

# 1

1. Green Valley Mills produces carpet at plants in St. Louis and Richmond. The carpet is then shipped to two outlets located in Chicago and Atlanta. The cost per ton of shipping carpet from each of the two plants to the two warehouses is as follows.

From	To	
	Chicago	Atlanta
St. Louis	\$40	\$65
Richmond	70	30



## TRANSPORTATION AND ASSIGNMENT SOLUTION METHODS

The plant at St. Louis can supply 250 tons of carpet per week; the plant at Richmond can supply 400 tons per week. The Chicago outlet has a demand of 300 tons per week, and the outlet at Atlanta demands 350 tons per week. The company wants to know the number of tons of carpet to ship from each plant to each outlet in order to minimize the total shipping cost. Solve this transportation problem.

5. Given a transportation problem with the following costs, supply, and demand, find the initial solution using the minimum cell cost method and Vogel's approximation model. Is the VAM solution optimal?

From	To			Supply
	1	2	3	
A	\$ 6	\$ 7	\$ 4	100
B	5	3	6	180
C	8	5	7	200
Demand	135	175	170	

8. Consider the following transportation problem.

From	To			Supply
	1	2	3	
A	\$ 6	\$ 9	\$ 7	130
B	12	3	5	70
C	4	11	11	100
<b>Demand</b>	80	110	60	

- Find the initial solution using the minimum cell cost method.
- Solve using the stepping-stone method.



13. A manufacturing firm produces diesel engines in four cities—Phoenix, Seattle, St. Louis, and Detroit. The company is able to produce the following numbers of engines per month.

<i>Plant</i>	<i>Production</i>
1. Phoenix	5
2. Seattle	25
3. St. Louis	20
4. Detroit	25

Three trucking firms purchase the following numbers of engines for their plants in three cities.

<i>Firm</i>	<i>Demand</i>
A. Greensboro	10
B. Charlotte	20
C. Louisville	15

The transportation costs per engine (\$100s) from sources to destinations are shown in the following table. However, the Charlotte firm will not accept engines made in Seattle, and the Louisville firm will not accept engines from Detroit; therefore, these routes are prohibited.

<b>From</b>	<b>To</b>		
	<i>A</i>	<i>B</i>	<i>C</i>
1	\$ 7	\$ 8	\$ 5
2	6	10	6
3	10	4	5
4	3	9	11

- Set up the transportation tableau for this problem. Find the initial solution using VAM.
- Solve for the optimal solution using the stepping-stone method. Compute the total minimum cost.
- Formulate this problem as a linear programming model.

21. A large manufacturing company is closing three of its existing plants and intends to transfer some of its more skilled employees to three plants that will remain open. The number of employees available for transfer from each closing plant is as follows.

Closing Plant	Transferable Employees
1	60
2	105
3	<u>70</u>
Total	235

The following number of employees can be accommodated at the three plants remaining open.

Open Plants	Employees Demanded
A	45
B	90
C	<u>35</u>
Total	170

Each transferred employee will increase product output per day at each plant as shown in the following table. The company wants to transfer employees so as to ensure the maximum increase in product output.

From	To		
	<i>A</i>	<i>B</i>	<i>C</i>
1	5	8	6
2	10	9	12
3	7	6	8

- Find the initial solution using VAM.
- Solve using MODI.

34. Orient Express is a global distribution company that transports its clients' products to customers in Hong Kong, Singapore, and Taipei. All of the products Orient Express ships are stored at three distribution centers, one in Los Angeles, one in Savannah, and one in Galveston. For the coming month the company has 450 containers of computer components available at the Los Angeles center, 600 containers available at Savannah, and 350 containers available in Galveston. The company has orders for 600 containers from Hong Kong, 500 containers from Singapore, and 500 containers from Taipei. The shipping costs per container from each U.S. port to each of the overseas ports are shown in the following table.

U.S. Center Distribution	Overseas Port		
	<i>Hong Kong</i>	<i>Singapore</i>	<i>Taipei</i>
Los Angeles	\$300	\$210	\$340
Savannah	490	520	610
Galveston	360	320	500

The Orient Express as the overseas broker for its U.S. customers is responsible for unfulfilled orders, and it incurs stiff penalty costs from overseas customers if it does not meet an order. The Hong Kong customers charge a penalty cost of \$800 per container for unfulfilled demand, Singapore customers charge a penalty cost of \$920 per container, and Taipei customers charge \$1,100 per container. Formulate and solve a transportation model to determine the shipments from each U.S. distribution center to each overseas port that will minimize shipping costs. Indicate what portion of the total cost is a result of penalties.